

Quantum Channels (Part 1)

Recap:

Quantum (Information) Theory in a slide....

Given a quantum system S , a quantum state is any Hermitian matrix ρ such that:

1. (Non-negativity) $\text{eigs}(\rho) \geq 0$.
2. (Normalization) $\text{tr}(\rho) = 1$.

Quantum states

Quantum evolutions

Given a quantum system S in a state ρ any quantum evolution is described by a linear transformation $\rho \rightarrow \mathcal{E}(\rho)$, such that $\sigma = \mathcal{E}(\rho)$ is always another valid quantum state.

last part
of the
note
so far

Given a quantum system S of dimension d , any general quantum measurement \mathcal{M} on S with m outcomes is described by a set of matrices $\mathcal{M} = \{M_0, M_1, \dots, M_{m-1}\}$ such that

1. (Non-negativity) $\text{eigs}(M_i) \geq 0$ for all $i = 0, 1, \dots, m-1$.
2. (Normalization) $\sum_i M_i = \mathbb{1}$.

Quantum measurements

Where $\mathbb{1}$ is the $d \times d$ identity matrix. Moreover, given a quantum state ρ , the probability of the k 'th outcome of doing the measurement \mathcal{M} on the state is given by

$$\text{Prob}[k] = \text{tr}[M_k \rho]. \quad (16)$$

Basic requirements on \mathcal{E} & the Kraus representation

① **Linearity**: $\mathcal{E}(\rho\rho + (1-\rho)\sigma) = \rho\mathcal{E}(\rho) + (1-\rho)\mathcal{E}(\sigma)$

Intuition ① Toss a coin & prepare ρ or σ and then apply \mathcal{E} . Output resulting state.

② Apply \mathcal{E} to ρ and σ , toss a coin and then output $\mathcal{E}(\rho)$ or $\mathcal{E}(\sigma)$.

Linearity captures the fact that the state output from procedures ① & ② are entirely indistinguishable.

2) Output state needs to be a real state

- i. $\text{Tr}(\mathcal{E}(\rho)) = 1$ (normalized)
- ii. $\mathcal{E}(\rho) = \mathcal{E}(\rho)^\dagger$ (Hermitian)
- iii. $\text{eig}(\mathcal{E}(\rho)) \geq 0$ (positivity / non-negativity)

These requirements on \mathcal{E} are satisfied if

$$\mathcal{E}(\rho) = \sum_i A_i \rho A_i^\dagger$$

$$\text{s.t. } \sum_i A_i^\dagger A_i = I$$

In just this 'if' can be upgraded to an 'only if' if we require that the output of the channel applied to just a subsystem in a real state. Will come back to this later.

This is called the Kraus Representation

Proof

$$\begin{aligned} \mathcal{E}(\rho \rho + (1-\rho)\sigma) &= \sum_i A_i (\rho \rho + (1-\rho)\sigma) A_i^\dagger \\ &= \rho \sum_i A_i \rho A_i^\dagger + (1-\rho) \sum_i A_i \sigma A_i^\dagger \\ &= \rho \mathcal{E}(\rho) + (1-\rho) \mathcal{E}(\sigma) \end{aligned} \quad \checkmark$$

$$2) \text{i. } \text{Tr}(\mathcal{E}(\rho)) = \text{Tr}(\sum_i A_i \rho A_i^\dagger)$$

$$= \sum_i \text{Tr}(A_i \rho A_i^\dagger)$$

$$= \sum_i \text{Tr}(A_i^\dagger A_i \rho)$$

$$= \text{Tr}(\sum_i A_i^\dagger A_i \rho)$$

$$= \text{Tr}(\rho)$$

trace preserved so if

$\text{Tr}(\rho) = 1$ then

$\text{Tr}(\mathcal{E}(\rho)) = 1$

2) ii. $\text{eig}(\sigma) \geq 0 \Leftrightarrow \langle \phi | \sigma | \phi \rangle \geq 0 \forall |\phi\rangle$

Any normalized
quantum state

$$\left(\begin{array}{l} \text{to see this write } \sigma = \sum_{\lambda} \lambda | \lambda \rangle \langle \lambda | \\ \langle \phi | \sigma | \phi \rangle = \sum_{\lambda} \lambda | \langle \phi | \lambda \rangle |^2 \\ \geq 0 \text{ if } \lambda \geq 0 \end{array} \right)$$

Showing that $\langle \phi | \mathcal{E}(\rho) | \phi \rangle \geq 0 \forall |\phi\rangle$

$$\begin{aligned} \langle \phi | \sum_i A_i \rho A_i^\dagger | \phi \rangle &= \sum_i \langle \phi | A_i \rho \overset{\curvearrowleft}{A_i^\dagger} | \phi \rangle \\ &= \sum_i \lambda \underbrace{|\langle \phi | A_i | \lambda \rangle|}_{\text{+ve}} \underbrace{|\langle \lambda | A_i^\dagger | \phi \rangle|}_{\text{+ve}} \\ &= \sum_i \lambda |\langle \phi | A_i | \lambda \rangle|^2 \\ &\therefore \text{+ve } \forall \rho \text{ +ve} \checkmark \end{aligned}$$

Examples of Common Channels

1) Unitary Dynamics $\mathcal{E}(\cdot) = \underset{\mathbb{C}}{\mathcal{U}}(\cdot) \mathcal{U}^\dagger$

$$\text{Check: } \sum_i A_i^\dagger A_i = \mathcal{U}^\dagger \mathcal{U} = \mathcal{I} \quad \mathcal{U} \mathcal{U}^\dagger = \mathcal{V}^\dagger \mathcal{V} = \mathcal{I}$$

(covered previously)
It's helpful to know how a unitary on a single qubit
affects a state on a Bloch sphere.

$$\text{Let } \rho = \frac{1}{2} (\mathcal{I} + \vec{\sigma} \cdot \vec{\Omega})$$

$$E(\rho) = U \rho U^\dagger = \underbrace{i(I + \underline{S} \cdot \underline{\sigma})}_{\underline{\Sigma}} = U(\underline{\Gamma} \cdot \underline{\sigma})U^\dagger$$

Claim: $\underline{S} = O \underline{\Gamma}$ where O is an orthogonal matrix

Orthogonal matrices induce (length preserving) rotations on real vectors. $\left\{ \begin{array}{l} \text{i.e. } O O^\dagger = O^\dagger O = I \\ \text{Real}(O) = O \end{array} \right.$

(i.e. they are the real analogue of unitaries)

Proof:

To show that O is orthogonal we just need to show that $|\underline{\Sigma}| = |\underline{\Gamma}|$

$$\begin{aligned} \text{To do so, we note that } \text{Tr}((\underline{S} \cdot \underline{\sigma})^2) &= \text{Tr}\left(\left(\sum_i s_i \sigma_i\right)^2\right) \\ &= \text{Tr}\left(\sum_{ij} s_i s_j \sigma_i \sigma_j\right) \\ &= \sum_{ij} s_i s_j \delta_{ij} \\ &= \sum_i s_i^2 = |\underline{S}|^2 \end{aligned}$$

Now we note that

$$\begin{aligned} |\underline{\Sigma}|^2 &= \text{Tr}((\underline{S} \cdot \underline{\sigma})^2) = \text{Tr}((U(\underline{\Gamma} \cdot \underline{\sigma})U^\dagger)^2) = \text{Tr}(U(\underline{\Gamma} \cdot \underline{\sigma})^2 U^\dagger) = \sum_{ij} \text{Tr}(\Gamma_i \Gamma_j U \sigma_i \sigma_j U^\dagger) \\ &= \sum_i \Gamma_i^2 = |\underline{\Gamma}|^2 \end{aligned}$$

$\therefore |\underline{S}| = |\underline{\Gamma}| \Rightarrow O$ is orthogonal

To get a better handle on the nature of this rotation, note that :

Any unitary on a qubit can be written as

$$U = e^{i\theta \mathbf{n} \cdot \boldsymbol{\sigma}}$$

(if this isn't immediately obvious, note that any unitary can be written as

$$U = e^{-iHt}$$

and H can always be expanded in the Pauli basis)

which is equivalent to

$$U = e^{i\theta \mathbf{n} \cdot \boldsymbol{\sigma}} = \cos(\theta) \mathbf{I} + i \sin(\theta) \mathbf{n} \cdot \boldsymbol{\sigma}$$

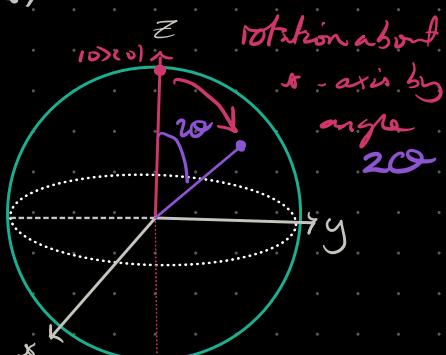
(if you've not shown this before, it's worth proving it to yourself)

What is the effect of evolving $|10\rangle\langle 01|$ under $e^{-i\theta \sigma_x}$?

$$\begin{aligned} e^{-i\theta \sigma_x/2} |10\rangle &= \cos(\theta) |10\rangle + i \sin(\theta) \sigma_x |10\rangle \\ &= \cos(\theta) |10\rangle + i \sin(\theta) |11\rangle \end{aligned}$$

$$\text{eg. } e^{+\frac{\pi}{2} \sigma_x} |10\rangle = \cos\left(\frac{\pi}{4}\right) |10\rangle + i \sin\left(\frac{\pi}{4}\right) |11\rangle = |11\rangle$$

Exercise : Show that $e^{-i\theta \mathbf{n} \cdot \boldsymbol{\sigma}}$ corresponds to a rotation about \mathbf{n} vector by an angle 2θ



2) Isometric Dynamics $\mathcal{E}(\cdot) = U(\cdot)U^\dagger$

Check: $\sum_i A_i^\dagger A_i = U^\dagger U = \mathcal{I}$ Note we only use $U^\dagger U = \mathcal{I}$
 is not $UU^\dagger = \mathcal{I}$
 $\therefore U$ could be an isometry
 (that isn't also unitary)

$$\text{eg. consider } U|10\rangle = \frac{1}{\sqrt{2}}(|10\rangle + |12\rangle)$$

$$U|12\rangle = \frac{1}{\sqrt{2}}(|12\rangle + |13\rangle)$$

This maps the system from a 2d Hilbert space $\{|10\rangle, |12\rangle\}$
 to a 4d space $\{|10\rangle, |12\rangle, |12\rangle, |13\rangle\}$

3) Convex combination of Unitaries

$$\mathcal{E}(\cdot) = \sum_i p_i U_i(\cdot) U_i^\dagger$$

$$0 \leq p_i \leq 1 \quad U_i^\dagger U_i = U_i^\dagger U_i = \mathcal{I}$$

Read: "apply unitary U_i with probability p_i "

$$A_i = \sqrt{p_i} U_i$$

$$\text{Check: } \sum_i A_i^\dagger A_i = \sum_i p_i U_i^\dagger U_i = \sum_i p_i \mathcal{I} = \mathcal{I} \quad \checkmark$$

4) ^{completely} Dephasing:

$$A_0 = |0\rangle\langle 0|$$

$$A_1 = |1\rangle\langle 1|$$

$$\text{Check } \sum_i A_i^\dagger A_i = |1\rangle\langle 1| + |0\rangle\langle 0| = \mathcal{I} \quad \checkmark$$

What does this channel do?

$$\rho = \frac{1}{2} (I + \Sigma \cdot \sigma) = \frac{1}{2} \begin{pmatrix} x+z & x-iy \\ x+iy & x-z \end{pmatrix}$$

$$\begin{aligned} E(\rho) &= \langle 0 | \rho | 0 \rangle | 0 \rangle \langle 0 | + \langle 1 | \rho | 1 \rangle | 1 \rangle \langle 1 | \\ &= \frac{1}{2} \begin{pmatrix} x+z & 0 \\ 0 & x-z \end{pmatrix} \end{aligned}$$

\Rightarrow Dephasing map kills the off-diagonal components & returns a classical state.

5) Completely Depolarizing Map

Given some fixed state τ

$$E(\rho) = \tau \sqrt{\rho}$$

This is a valid channel because it's linear & outputs a legitimate state.

So it must be possible to write it in terms of Kraus operators...

To see how just write τ in terms of its eigendecomposition

$$\tau = \sum_i \lambda_i |x_i\rangle \langle x_i|$$

$$\text{Then we can use: } A_{ij} = \sqrt{\lambda_i} |x_i\rangle \langle j|$$

To see why this works note that

$$\begin{aligned} E(\rho) &= \sum_{ij} \lambda_i |x_i\rangle \langle x_i| \rho |x_i\rangle \langle x_i| \\ &= \text{Tr}(\rho) \tau = \tau \checkmark \end{aligned}$$

$$\text{Also check } \sum_{ij} A_{ij}^* A_{ij} = \sum_{ij} \langle x_i | x_j \rangle \langle x_j | x_i \rangle = \sum_i \langle x_i | \sum_j \langle x_j | x_j \rangle = \sum_i \langle x_i | I = I$$

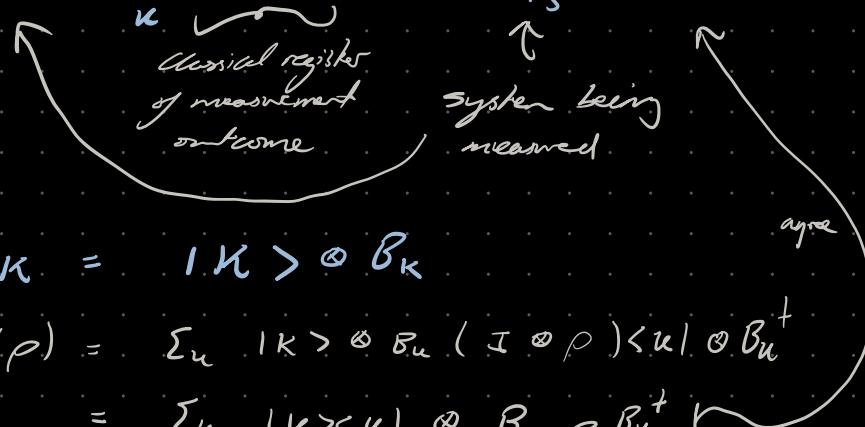
Exercise: What about a non-completely polarizing map?

e.g. one that outputs T with probability P
 $\otimes P \quad " " \quad 1-P$

What are the Kraus operators now?

6) The "Measure-record-and-update" channel

$$\mathcal{E}(\rho) = \sum_k |k\rangle \langle k|_x \otimes B_k \rho B_k^*$$



$$A_k = |k\rangle \langle k|_x \otimes B_k$$

$$\mathcal{E}(\rho) = \sum_k |k\rangle \langle k|_x (I \otimes \rho) \langle k|_x \otimes B_k^*$$

$$= \sum_k |k\rangle \langle k|_x \otimes B_k \rho B_k^*$$

$$\text{Check: } \sum_k A_k^* A_k = \sum_k \underbrace{\langle k| k \rangle}_\text{if } \sum_k B_k^* B_k = I$$

i.e. $\{A_k\}$ is a legit set of Kraus operators iff $\{B_k\}$ are a set of Kraus operators.

What has this got to do with measurements?

Let's look at the reduced output state on the classical register

$$\begin{aligned} \text{Tr}_s(\mathcal{E}(\rho)) &= \text{Tr}_s \left(\sum_u |u\rangle\langle u| \otimes B_u \rho B_u^\dagger \right) \\ &= \sum_u |u\rangle\langle u| \text{Tr}_s \left(\underbrace{B_u \rho B_u^\dagger}_{M_u = B_u^\dagger B_u} \right) \quad \downarrow M_u = B_u^\dagger B_u \\ &= \sum_u \text{Tr}(M_u \rho) |u\rangle\langle u| \quad \downarrow P_u = \text{Tr}(M_u \rho) \\ &= \sum_u P_u |u\rangle\langle u| \end{aligned}$$

Because $M_u^\dagger = (B_u^\dagger B_u)^\dagger = B_u B_u^\dagger = M_u$ is Hermitian it has real eigenvalues. Moreover, these eigenvalues are necessarily non-negative (because $\langle \phi | B_u^\dagger B_u | \phi \rangle = \| B_u | \phi \rangle \|_2^2 \geq 0 \forall \phi$)

So M_u is a ~~the~~ Hermitian operator

$\Rightarrow \sum_u M_u = \sum_u B_u^\dagger B_u = \mathcal{I}$ by assumption

$\therefore \{M_u\}$ defines a POVM

$\& P_u = \text{Tr}(M_u \rho)$ is just the probability of getting the k_u outcome

If we measure the classical register, as it is in the state

$\sum_u P_u |u\rangle\langle u|$ we get the k_u outcome with probability P_u
i.e. the measurement outcome has been recorded in the classical register X .

What if we trace out the second register instead?

Then we get

$$\begin{aligned} \text{Tr}_X(E(\rho)) &= \text{Tr}_X \left(\sum_k |k\rangle\langle k|_X \otimes B_k \rho B_k^+ \right) \\ &= \sum_k B_k \rho B_k^+ \quad \left\{ \begin{array}{l} \text{form of a standard} \\ \text{channel} \end{array} \right. \end{aligned}$$

$$E(\rho) = \sum_k \underbrace{|k\rangle\langle k|_X}_{\substack{\text{classical register} \\ \text{of measurement} \\ \text{outcome}}} \otimes \underbrace{B_k \rho B_k^+}_{\substack{\uparrow \\ \text{measurement update} \\ \text{rate (classically correlated} \\ \text{with corresponding recorded} \\ \text{measurement outcome)}}}$$

That is, if we look at the state of the system conditioned on obtaining measurement outcome $|k\rangle\langle k|_X$ we have

$$\begin{aligned} \rho &\rightarrow E(\rho) \rightarrow \underbrace{|k\rangle\langle k|_X \otimes B_k \rho B_k^+}_{\substack{\text{"get } k\text{"} \\ \text{ie. } (|k\rangle\langle k| \otimes I) E(\rho) (|k\rangle\langle k| \otimes I)}} \end{aligned}$$

This is not a normalized state, the normalized output state would be

$$\rho_{\text{out}} = |k\rangle\langle k| \otimes \rho_k \quad \text{with} \quad \rho_k = \frac{B_k \rho B_k^+}{\text{Tr}(B_k \rho)}$$

i.e. a POVM $\{M_u\}$ such that $M_u = B_u^+ B_u$ tells us two things

- ① The probability of getting output k $\text{Tr}(M_k \rho)$
- ② An update rule

$$\rho \rightarrow \rho_x = \frac{B_x \rho B_x^+}{\text{Tr}(M_x \rho)}$$

But note this update rule is not unique.

Say $B_u \rightarrow \tilde{B}_u = \underbrace{U B_u}_{\text{unitary isometry}} \text{ some}$

$$\tilde{B}_u^+ \tilde{B}_u = B_u^+ U^+ U B_u = B_u^+ B_u = M_u$$

Thus it's important to draw a distinction between the POVM, the measurement, which is all about how to extract information from a state & the (optional) state update rule which concerns the evolution of the state.



7) Partial Trace $E(\rho_{AB}) = \text{Tr}_B(\rho_{AB}) = \rho_A$

$$A_u = I_A \otimes \langle u | \circ$$

Check it works $\sum_u A_u \rho A_u^+ = \sum_u (I_A \otimes \langle u |) \rho (I_A \otimes |u \rangle)$

$$:= \text{Tr}_B(\rho) \checkmark$$

$$\sum_k A_k^\dagger A_k = I_A \otimes \sum_u |u \rangle \langle u| = I_{AB} \checkmark$$

8) State Preparation

$$\xrightarrow{\text{trivial input}} \rho_{\text{state to be prepared}}$$

$$A_K = \sqrt{\lambda_i} | \lambda_i \rangle$$

$$\text{Check it works} \quad \sum_n A_n (\mathbb{I}) A_n^\dagger = \sum_n \lambda_n | \lambda_n \rangle \langle \lambda_n |$$

$$\text{Check} \quad \sum_n A_n^\dagger A_n = \sum_n \lambda_n \langle \lambda_n | \lambda_n \rangle = \mathbb{I}$$

Note :

Combinations of channels are channels

$$(\mathcal{F} \circ \mathcal{E})(\rho) = \mathcal{F}(\mathcal{E}(\rho)) = \mathcal{F}\left(\sum_i A_i \rho A_i^\dagger\right)$$

$$\sum_j B_j^\dagger \quad \sum_i A_i^\dagger \quad = \sum_{ij} B_j^\dagger A_i \rho A_i^\dagger B_j^\dagger$$

Can introduce new Kraus operators $\rightarrow \sum_{i,j} C_{ij}$
(for the concatenated channel)

$$= \sum_k C_k \rho C_k^\dagger$$

$$\text{Check:} \quad \sum_k C_k^\dagger C_k = \sum_{ij} A_i^\dagger B_j^\dagger B_j A_i = \sum_i A_i^\dagger \sum_j B_j^\dagger B_j A_i = \mathbb{I}$$

Punchline : Basically everything that happens in quantum theory can be thought of as a quantum channel !

- State preparation
- Evolution
- Discarding a system
- Measurement
- Updates after measurement
- Interactions with an environment

more on this
to come

⇒ This unifying perspective, as we will see, can at times be helpful.